



# **Gevallestudies Wiskundige Ingieurstechnieken**

**Joppe Bos**

March 2025





**Joppe W. Bos**

Cryptographic Researcher and  
Technical Director at NXP  
Semiconductors

Secretary of the IACR (2017-  
2019, 2020-2022)

Editor of the Cryptology ePrint  
Archive (2019-today)

Editor-in-Chief of the IACR  
Communications in Cryptology

♥ Leuven, Belgium

## WHOAMI

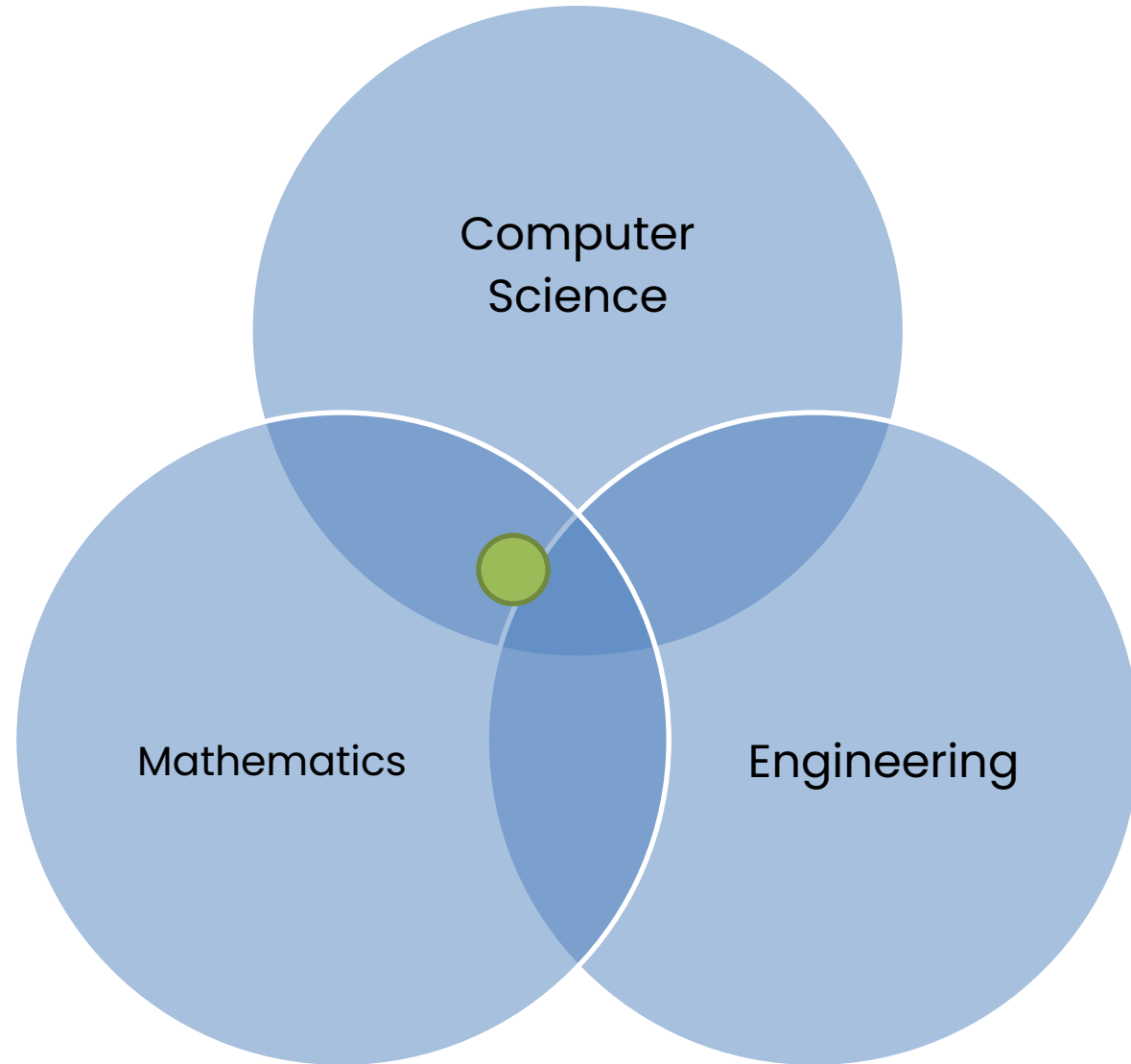
- Cryptographic researcher + Technical Director
  - in the competence center crypto & security at NXP Semiconductors, Leuven
  - Lead the PQC team
  - Lead security + crypto funded projects & university relations
- Post-doc
  - Cryptography Research Group at Microsoft Research, Redmond, USA.
- PhD in Cryptology
  - EPFL, Lausanne, Switzerland
- Bachelor / Master in Computer Science
  - University of Amsterdam



# Public Key Cryptography

Computational  
number theory

Number  
theoretic  
transform







# Breaking ECC

112-bit ECDLP  
solved using 224  
PlayStation 3  
game consoles.



## NXP Corporate Overview

Together we accelerate  
the **breakthroughs** that  
advance our world

We design purpose-built,  
rigorously tested technologies  
that enable devices to sense,  
think, connect and act  
intelligently to improve  
people's daily lives.



Automotive



Industrial & IoT



Mobile



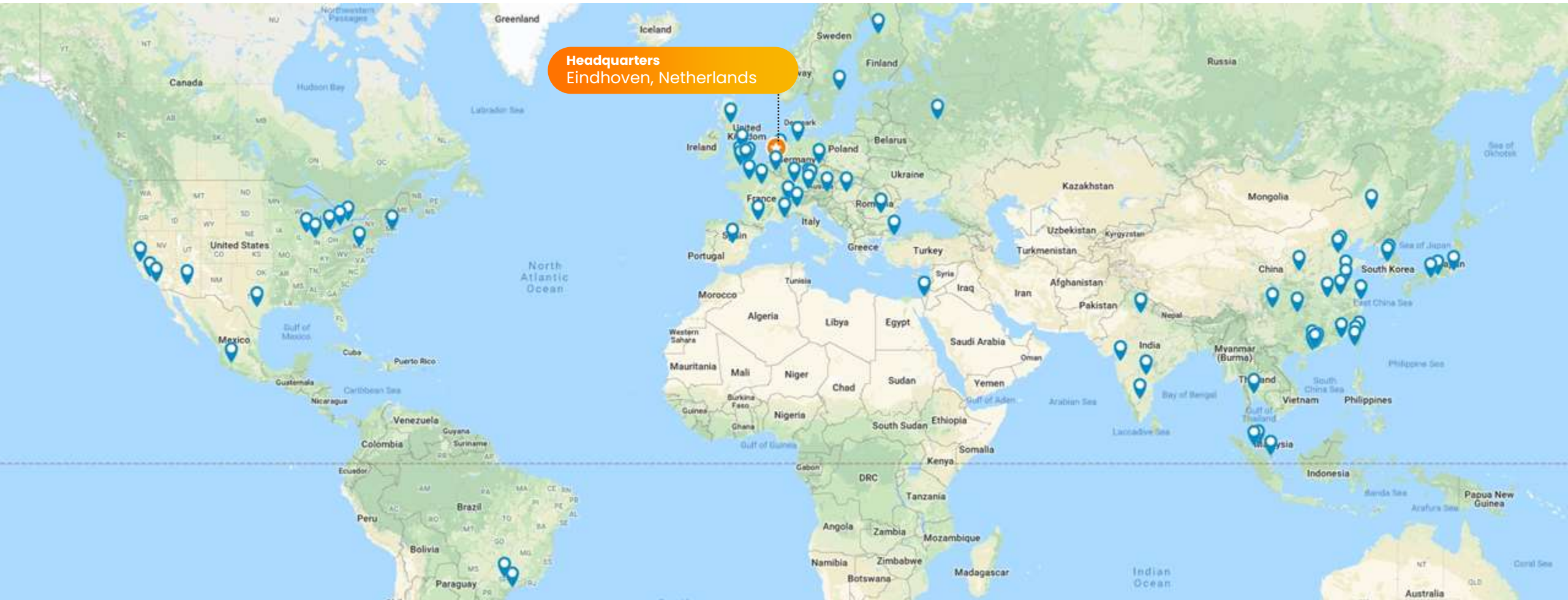
Communication  
Infrastructure





# NXP locations

~34,200 team members with operations in more than 30 countries





# Automotive market positions

Automotive

## Technology Leadership

+

## Applications Leadership

- #1 Auto processors
- #1 Auto applications processors
- #1 Auto RF
- #1 Auto DSPs
- #1 Cross-domain processors

- #1 Infotainment
- #1 Car radio
- #1 Secure car access
- #1 In-vehicle networking

Sources: Strategy Analytics: Automotive Semiconductors Vendor Market Shares, April 2024, Strategy Analytics: Infotainment and Telematics Semiconductors Vendor Market Shares, April 2024, Gartner: Semiconductors Market Shares, April 2024, S&P: competitive landscaping tool, April 2024, IHS: automotive semiconductors market tracker, April 2024





# Edge processing – a distributed intelligence pyramid

Millions

Cloud  
Data centers



**Public:** AWS, Azure, etc.  
**Private:** Salesforce, SAP, Honeywell, etc.  
**Hybrid:** public + private data centers

10's to 100's Millions

Network Edge  
Network computing



Billions

Application Edge  
IoT end points



Edge processing  
served market



# End-to-end solutions for Matter

A unified IP-based protocol to securely and robustly connect smart devices with each other, regardless of brand, and across smart home platforms

**Bring interoperability** in the Smart Home industry

**Simplify development** for “things”

**Increase reliability** for consumers

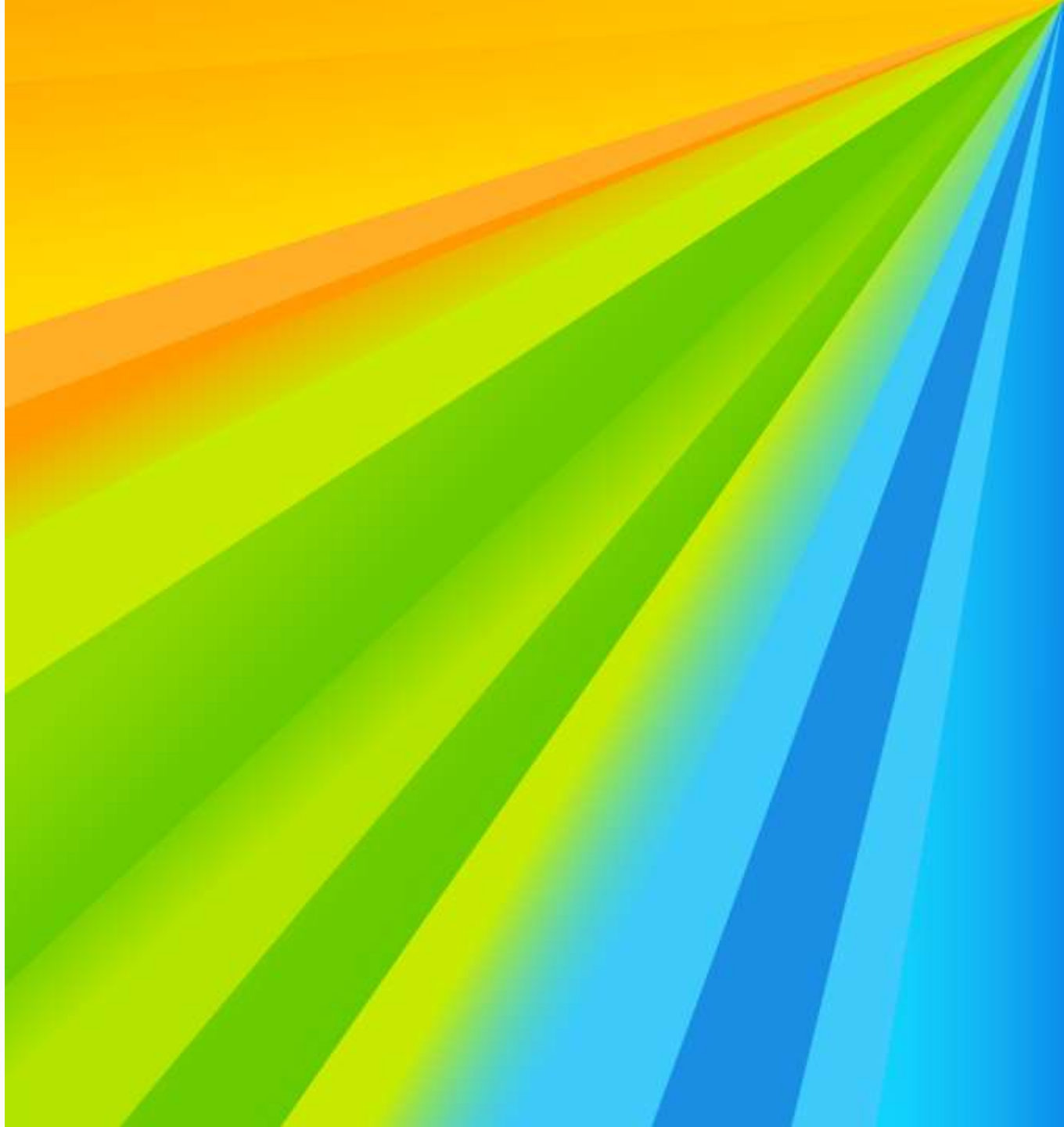
**Ensure security and privacy**

Led by global brands and 200+ companies





# Classical Cryptography





# Public-Key Cryptography

In **public-key** cryptography the theoretical foundation of the schemes used are problems which are **believed** to be hard

- Integer factorization problem (RSA)
- Discrete logarithm problem (DSA, ElGamal)

One of the main ingredients to these problems is a group

RSA  $\rightarrow (\mathbb{Z}/N\mathbb{Z})^\times \rightarrow$  integers  $[1, 2, \dots, N - 1]$  which are co-prime to  $N$

DSA/ElGamal  $\rightarrow \mathbb{F}_p^\times = (\mathbb{Z}/p\mathbb{Z})^\times \rightarrow$  integers  $[1, 2, \dots, p - 1]$  where  $p$  is prime

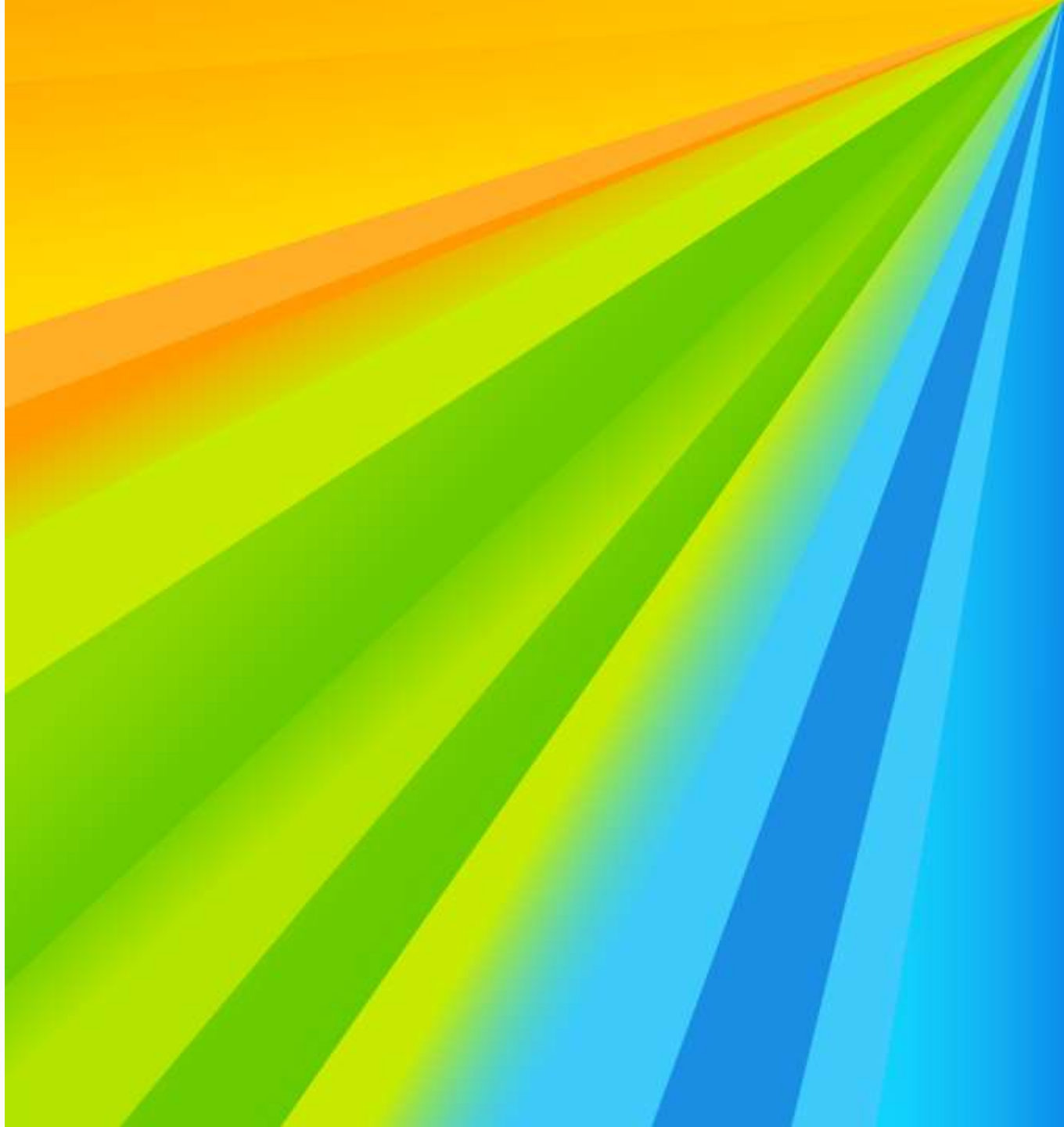
Elliptic Curve Cryptography  $\rightarrow E/\mathbb{F}_p \rightarrow$  point on  $E(\mathbb{F}_p)$  where  $p$  is prime



Application	Encryption Scheme, Signature Scheme, Identification Scheme, etc.		
Cryptosystem	DSA, ElGamal, Schnorr, etc.		RSA, Rabin, etc.
Computational Problem	The Discrete Logarithm Problem in a Group of prime Order		The Factoring Problem
Algebraic Structure	The multiplicative group of integers modulo a prime	Elliptic Curve Group over a Finite Field	The set of integers modulo the product of two primes



# Post-Quantum Cryptography





## How IBM's new five-qubit universal quantum computer works

IBM achieves an important milestone with new quantum computer in the cloud.

CHRIS LEE NEWS | 23 October 2019

### Hello quantum world! Google publishes landmark quantum supremacy claim

The company says that its quantum computer is the first to perform a calculation that would be practically impossible for a classical machine.

Elizabeth Gibney



### Eagle's quantum performance progress

Last November, IBM Quantum announced Eagle, a 127-qubit quantum processor based on the transmon superconducting qubit architecture. The IBM Quantum team adapted advanced semiconductor signal delivery and packaging into a technology node to develop superconducting quantum processors.

quantum  
new chip was  
performance.



## NXP, eleQtron and ParityQC Reveal their First Quantum Computing Demonstrator for the DLR Quantum Computing Initiative

May 30, 2024 2:00 PM CEST (UTC+2) by NXP Semiconductors Press Release

### Quantum error correction below the surface code threshold

Google Quantum AI and Collaborators  
(Dated: August 27, 2024)

SHARE



in



- NXP, eleQtron and ParityQC reveal their first quantum computing demonstrator for the DLR Quantum Computing Initiative
- It was commissioned by the DLR Quantum Computing Initiative (DLR QCI) to expand the quantum expertise of its partners from research and industry



# Security impact of quantum computers

## Requirements: Cryptography

Asymmetric	Symmetric
RSA-3072	AES-128
ECC P-256	SHA-256



"All use of cryptography must use an algorithm that meets at least 128 bits of security."





# Quantum potential to destroy security as we know it



## **Confidential email messages, private documents, and financial transactions**

Secure today but could be compromised in the future, even if encrypted



## **Firmware update mechanisms in vehicles**

Could be circumvented and allow dangerous modifications



## **Critical industrial and public service infrastructure (for healthcare, utilities, and transportation using internet and virtual private networks)**

Could become exposed – potentially destabilize cities



## **Audit trails and digitally signed documents associated with safety (auto certification and pharmaceutical authorizations)**

Could be retrospectively modified



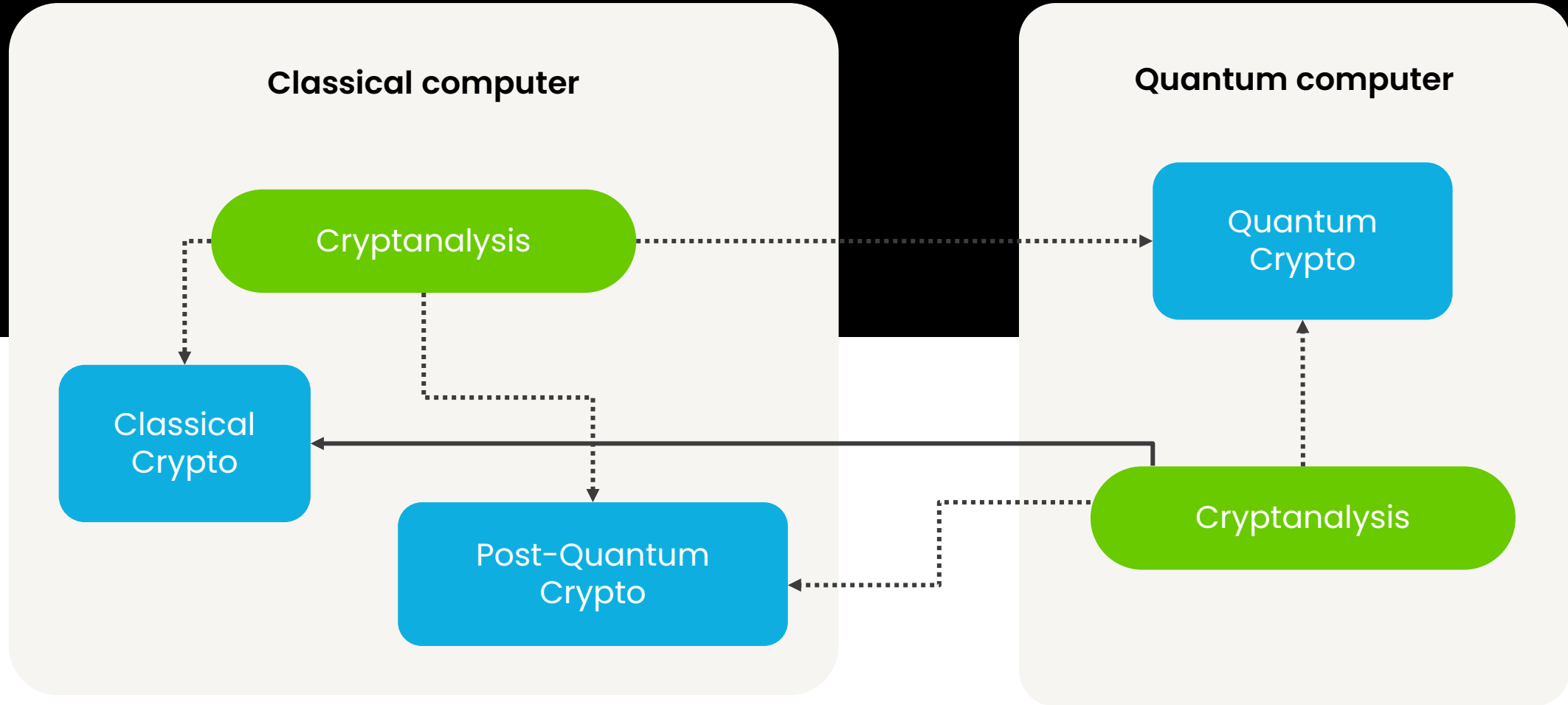
## **The integrity of blockchains**

Could be retrospectively compromised – could include fraudulent manipulation of ledger and cryptocurrency transactions





# Post-quantum versus quantum crypto







## Post-Quantum Cryptography

### Requirement 1

Run on  
classical hardware

### Requirement 2

Be secure against adversaries  
armed with classical computers

### Requirement 3 NEW

Be secure against adversaries  
armed with quantum computers

### Requirement 4

Be secure against Side-Channel Analysis (SCA)  
and Fault Injection (FI) attacks



# Is Post-Quantum Cryptography relevant for you?

## Standards & Compliance

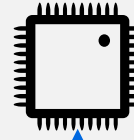


NIST



## Crypto Agility

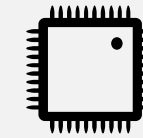
PQC RoT



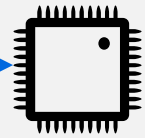
Secure Updates



## Store Now Decrypt Later

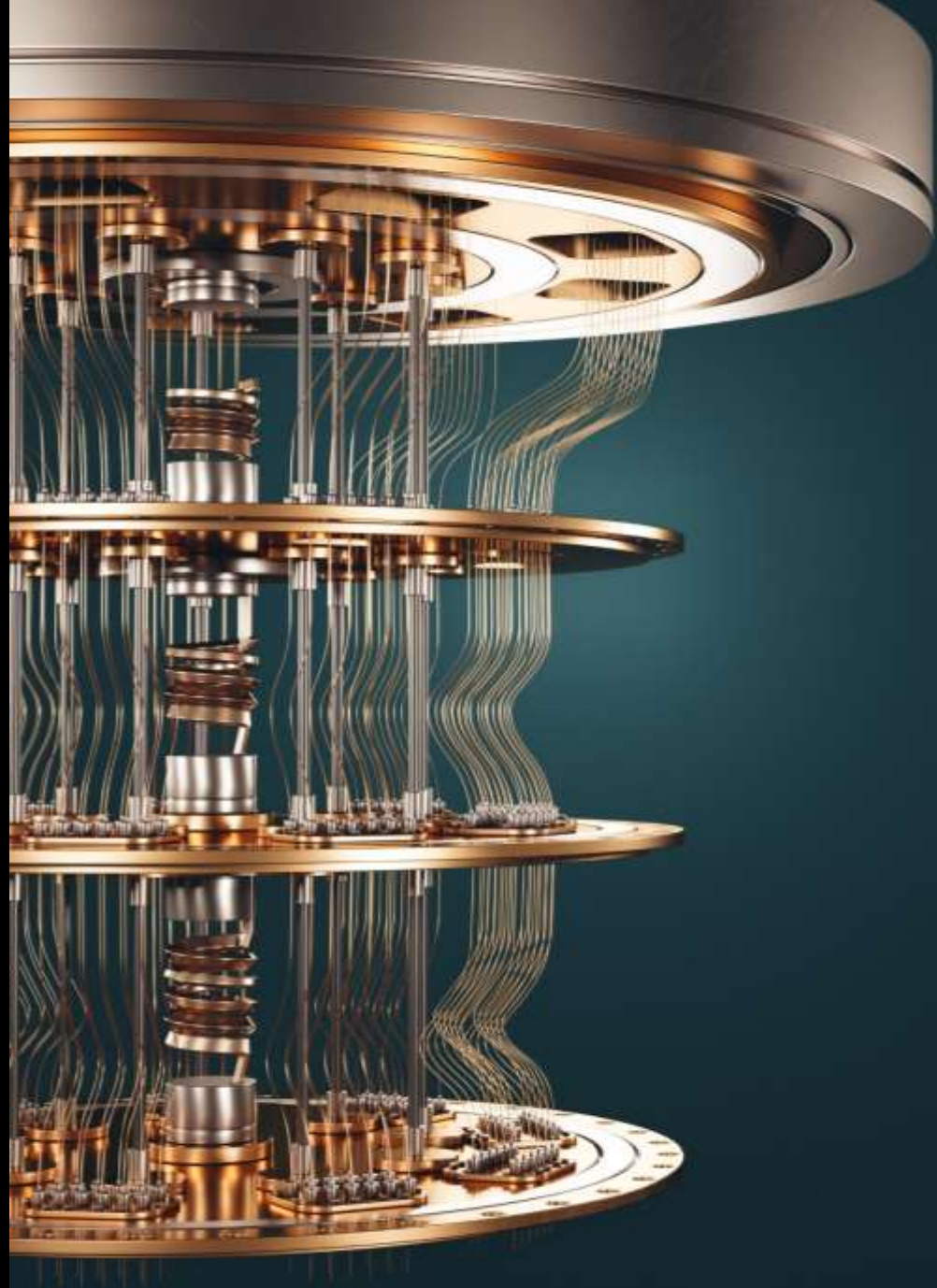


TLS 1.3





**Post-quantum  
crypto standards  
are coming  
It doesn't matter if  
you believe in  
quantum  
computers or not**

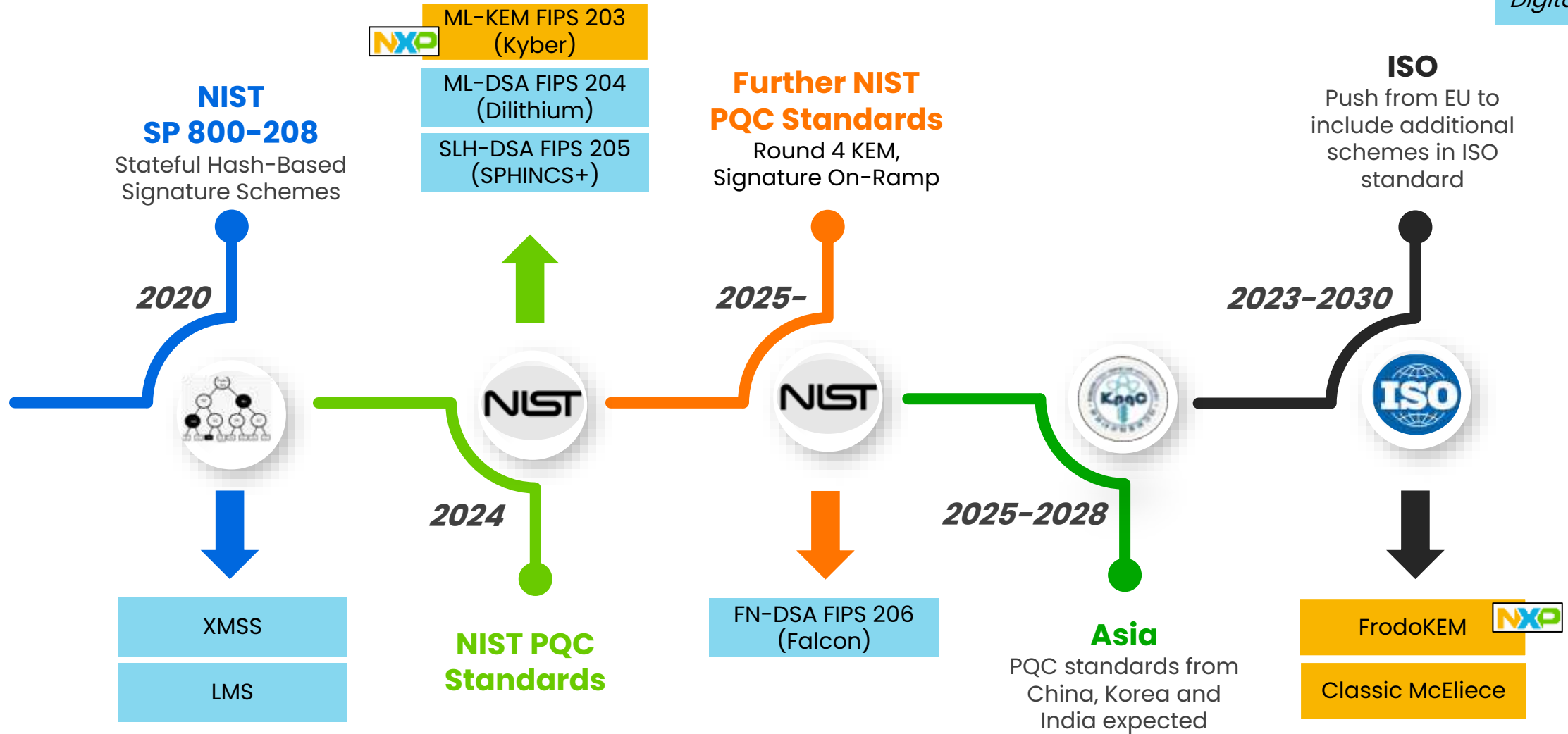




# PQC Algorithm Standardization

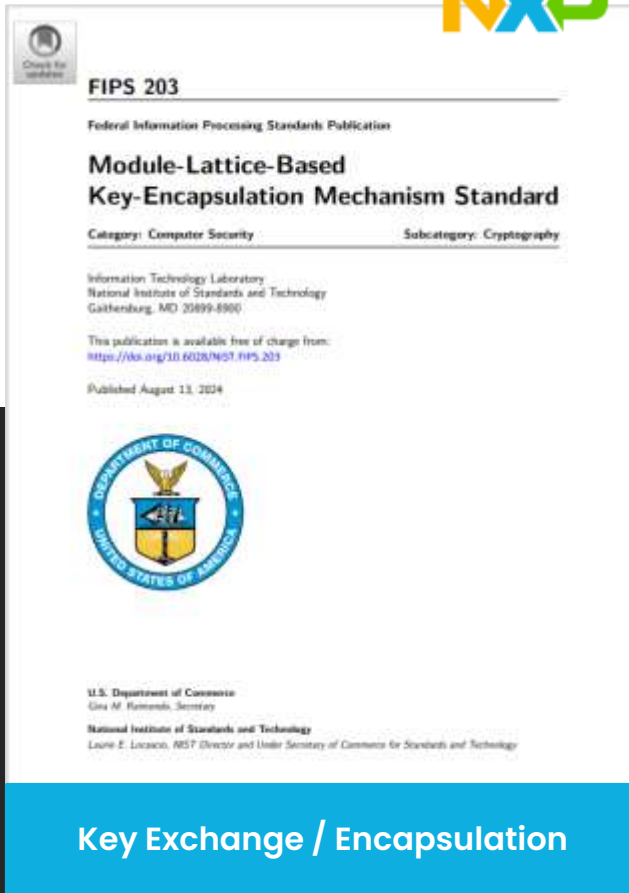
Key Exchange

Digital Signature





# New algorithms and standards

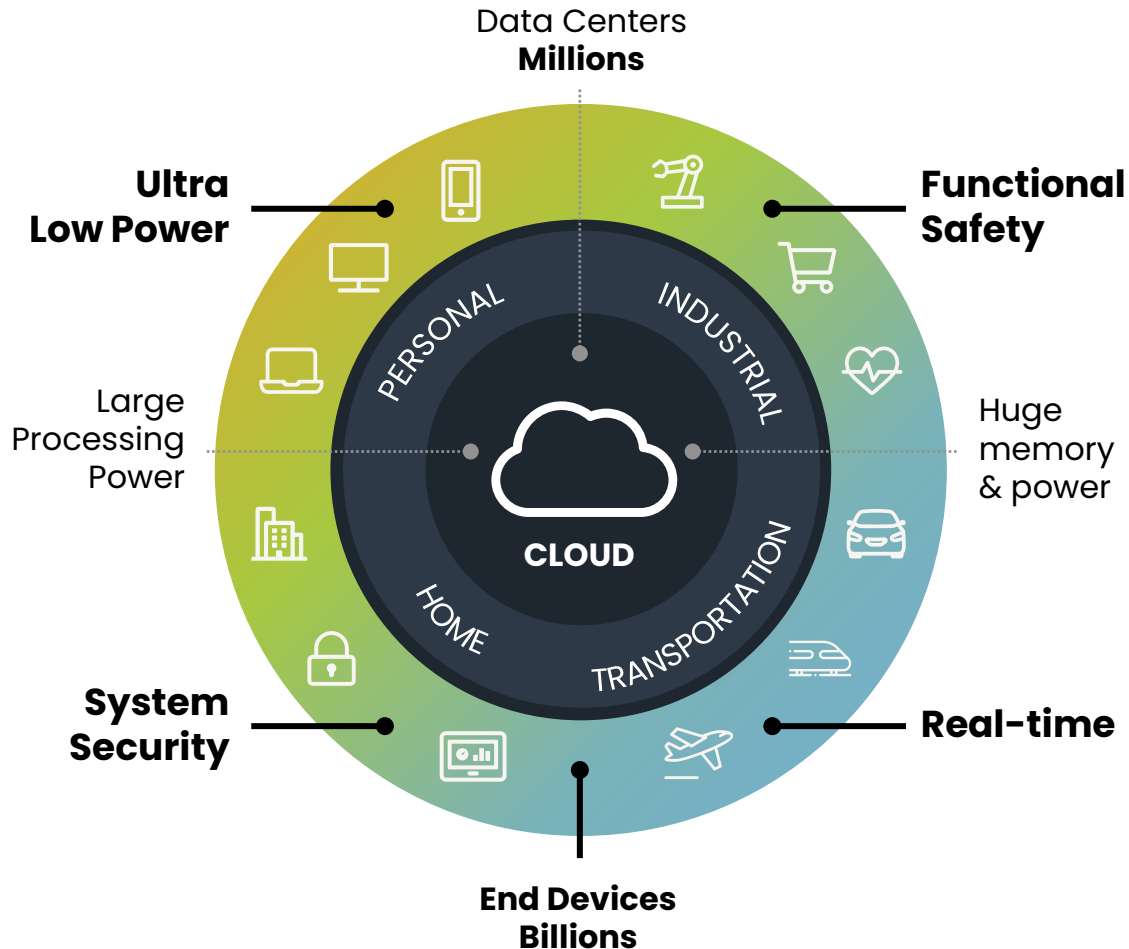


More ongoing and upcoming! FIPS 206, Round 4, On-Ramp, ISO, etc..

- [1] ML-KEM, <https://nvlpubs.nist.gov/nistpubs/fips/nist.fips.203.pdf>
- [2] ML-DSA, <https://nvlpubs.nist.gov/nistpubs/fips/nist.fips.204.pdf>
- [3] SLH-DSA, <https://nvlpubs.nist.gov/nistpubs/fips/nist.fips.205.pdf>
- [4] LMS / XMSS, <https://nvlpubs.nist.gov/nistpubs/SpecialPublications/NIST.SP.800-208.pdf>



# Impact PQC on our eco-system



Data collection, processing and decisions at the edge  
Devices securely connected to the cloud

## No Silver Bullet

If a crypto scheme was better, we would have standardized this already

## Cryptographic Keys

Orders of magnitude larger.

In the final: up to 1.3MB

Winners: up to 4.8KB

(ECC: 32 bytes, RSA: 384 bytes)

## Performance

Varies: some faster some significantly slower.

SHA-3 is a dominating component (~80%)

## Memory

Orders of magnitude more:

up 100KB memory of RAM when executing

NXP has dedicated implementations reaching ~16KB of RAM

## Bandwidth & Power

Larger signatures (up to 4.6KB)

→ more bandwidth required

→ increase in power usage



# Technical aspects of new algorithms

See pqm4 open source project for benchmarks! [A]  
Assuming Cortex-M4 @ 200 MHz software-only.  
For LMS numbers taken from Campos et al. [B]

Algorithm	PQC	Encaps	Decaps	SK	PK	CT	Algorithm
EC-P384	No	"Fast"	"Fast"	48 B	48 B	96 B	EC-P384
FIPS 203 (ML-KEM)	Yes	4 ms	4 ms	2 400 B	1 184 B	1 088 B	FIPS 203 (ML-KEM)

Algorithm	PQC	Encaps	Decaps	SK	PK	CT	Algorithm
ECDSA-P384	No	"Fast"	"Fast"	48 B	48 B	96 B	ECDSA-P384
FIPS 204 (ML-DSA)	Yes	31 ms	12 ms	4 032 B	1 952 B	3 309 B	FIPS 204 (ML-DSA)
FIPS 205 (SLH-DSA)***	Yes	77 s	68 ms	96 B	48 B	16 224 B	FIPS 205 (SLH-DSA)***
SP 800-20 (LMS/XMSS)	Yes	** (Stateful) 19 s	13 ms	48 B	48 B	1 860 B	SP 800-208 (LMS/XMSS)

\* NIST Level 3 parameter sets

\*\* Significant reduction possible by increasing memory consumption for state

\*\*\* New parameter sets coming that will improve performance & signature size!



# What is the impact on the billions of embedded devices?



Automotive

70%

70% connected cars by 2025



Industrial & IoT

12B

IoT Edge & end nodes from **6B units** in 2021 to **12B units** in 2025



Mobile

60B

Tagging **60B products** per year by 2025



Communication Infrastructure

40B

Secure anchors & services for **40B processors**



Automotive



eGovernment



Bank cards



Smart mobility (MIFARE) cards



Tags & Authentication



Readers



Mobile



# Cryptographic Suite for Algebraic Lattices (CRYSTALS)

The Cryptographic Suite for Algebraic Lattices (CRYSTALS) encompasses

- **Kyber**, a Key Encapsulation Mechanism (KEM) -> referred to in FIPS 203 as **ML-KEM**
- **Dilithium**, for Digital Signatures -> referred to in FIPS 204 as **ML-DSA**

Theory: same building blocks

- Module Learning with Errors
- Number-Theoretic Transformations

Many new techniques to deal with!

Kyber uses the 'Fujisaki-Okamoto Transform' to get strong security

Dilithium uses 'Rejection Sampling' as a core component for producing signatures





# Solving systems of linear equations

$$\begin{matrix} \mathbb{Z}_{13}^{7 \times 4} \\ \begin{array}{|c|c|c|c|} \hline 4 & 1 & 11 & 10 \\ \hline 5 & 5 & 9 & 5 \\ \hline 3 & 9 & 0 & 10 \\ \hline 1 & 3 & 3 & 2 \\ \hline 12 & 7 & 3 & 4 \\ \hline 6 & 5 & 11 & 4 \\ \hline 3 & 3 & 5 & 0 \\ \hline \end{array} \end{matrix} \times \begin{matrix} \text{secret} \\ \mathbb{Z}_{13}^{4 \times 1} \\ \begin{array}{|c|} \hline \\ \hline \\ \hline \\ \hline \end{array} \end{matrix} = \begin{matrix} \mathbb{Z}_{13}^{7 \times 1} \\ \begin{array}{|c|} \hline 4 \\ \hline 8 \\ \hline 1 \\ \hline 10 \\ \hline 4 \\ \hline 12 \\ \hline 9 \\ \hline \end{array} \end{matrix}$$

Linear system problem: given **blue**, find **red**



# Solving systems of linear equations

$$\begin{matrix} \mathbb{Z}_{13}^{7 \times 4} \\ \begin{array}{|c|c|c|c|} \hline 4 & 1 & 11 & 10 \\ \hline 5 & 5 & 9 & 5 \\ \hline 3 & 9 & 0 & 10 \\ \hline 1 & 3 & 3 & 2 \\ \hline 12 & 7 & 3 & 4 \\ \hline 6 & 5 & 11 & 4 \\ \hline 3 & 3 & 5 & 0 \\ \hline \end{array} \end{matrix} \times \begin{matrix} \text{secret} \\ \mathbb{Z}_{13}^{4 \times 1} \\ \begin{array}{|c|} \hline 6 \\ \hline 9 \\ \hline 11 \\ \hline 11 \\ \hline \end{array} \end{matrix} = \begin{matrix} \mathbb{Z}_{13}^{7 \times 1} \\ \begin{array}{|c|} \hline 4 \\ \hline 8 \\ \hline 1 \\ \hline 10 \\ \hline 4 \\ \hline 12 \\ \hline 9 \\ \hline \end{array} \end{matrix}$$

Easily solved using  
Gaussian elimination  
(Linear Algebra 101)

Linear system problem: given **blue**, find **red**



# Learning with errors problem

random

 $\mathbb{Z}_{13}^{7 \times 4}$

4	1	11	10
5	5	9	5
3	9	0	10
1	3	3	2
12	7	3	4
6	5	11	4
3	3	5	0

secret

 $\mathbb{Z}_{13}^{4 \times 1}$

6
9
11
11

small noise

 $\mathbb{Z}_{13}^{7 \times 1}$

0
-1
1
1
1
0
-1

$\mathbb{Z}_{13}^{7 \times 1}$

4
7
2
11
5
12
8



# Learning with errors problem

random  
 $\mathbb{Z}_{13}^{7 \times 4}$

4	1	11	10
5	5	9	5
3	9	0	10
1	3	3	2
12	7	3	4
6	5	11	4
3	3	5	0

$\times$

secret  
 $\mathbb{Z}_{13}^{4 \times 1}$


$+$

small noise  
 $\mathbb{Z}_{13}^{7 \times 1}$


$=$

$\mathbb{Z}_{13}^{7 \times 1}$

4
7
2
11
5
12
8

Computational LWE problem: given **blue**, find **red**



# Toy example versus real-world example

$$\mathbb{Z}_{13}^{7 \times 4}$$

4	1	11	10
5	5	9	5
3	9	0	10
1	3	3	2
12	7	3	4
6	5	11	4
3	3	5	0

$$\mathbb{Z}_{2^{15}}^{752 \times 8}$$

8			
...			
2738	3842	3345	2979
2896	595	3607	
377	1575		
2760			
...			

752

$$752 \times 8 \times 15 \text{ bits} = 11 \text{ KiB}$$



# Ring learning with errors problem

random  
 $\mathbb{Z}_{13}^{7 \times 4}$

4	1	11	10
10	4	1	11
11	10	4	1
1	11	10	4
4	1	11	10
10	4	1	11
11	10	4	1

Each row is the cyclic shift of the row above



# Ring learning with errors problem

random  
 $\mathbb{Z}_{13}^{7 \times 4}$

4	1	11	10
3	4	1	11
2	3	4	1
12	2	3	4
9	12	2	3
10	9	12	2
11	10	9	12

Each row is the cyclic  
shift of the row above

...

with a special wrapping rule:  
 $x$  wraps to  $-x \bmod 13$ .



# Ring learning with errors problem

random  
 $\mathbb{Z}_{13}^{7 \times 4}$

4	1	11	10
---	---	----	----

Each row is the cyclic  
shift of the row above

...

with a special wrapping rule:

$x$  wraps to  $-x \bmod 13$  ( $\rightarrow \mathbb{Z}_{13}[x]/\langle x^4 + 1 \rangle$ )

So I only need to tell you the first row.



# Ring learning with errors problem

$$\mathbb{Z}_{13}[x]/\langle x^4 + 1 \rangle$$

$$4 + 1x + 11x^2 + 10x^3$$

random

×

$$6 + 9x + 11x^2 + 11x^3$$

secret

+

$$0 - 1x + 1x^2 + 1x^3$$

small noise

---

=

$$10 + 5x + 10x^2 + 7x^3$$



# Ring learning with errors problem

$$\mathbb{Z}_{13}[x]/\langle x^4 + 1 \rangle$$

$$4 + 1x + 11x^2 + 10x^3$$

random

×

$$\text{secret}$$

secret

+

$$\text{small noise}$$

small noise

=

$$10 + 5x + 10x^2 + 7x^3$$

Computational ring-LWE problem: given **blue**, find **red**



# Basic ring-LWE-DH key agreement

- Reformulation of Peikert's ring-LWE KEM (*PQCrypto 2014*)

public: "big"  $a$  in  $R_q = \mathbb{Z}_q[x]/(x^n+1)$

**Alice**

secret:

random "small"  $s, e$  in  $R_q$

**Bob**


secret:

random "small"  $s', e'$  in  $R_q$

$$b = a \cdot s + e$$


$$b' = a \cdot s' + e'$$


shared secret:

$$s \cdot b' = s \cdot (a \cdot s' + e') \approx s \cdot a \cdot s'$$


shared secret:

$$b \cdot s' \approx s \cdot a \cdot s'$$


**These are only approximately equal  $\Rightarrow$  need rounding**



# What is the impact of PQC on Industrial IoT?





# From theory to practice: small-memory implementations

Do these implementations actually run on embedded systems?

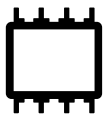
		pqm4	
		Runtime	RAM
Dilithium-2	Sign	19 ms	50 kB
	Verify	7 ms	11 kB
Dilithium-3	Sign	31 ms	69 kB
	Verify	12 ms	10 kB
Dilithium-5	Sign	42 ms	123 kB
	Verify	21 ms	12 kB



# From theory to practice: small-memory implementations

Do these implementations actually run on embedded systems?

		pqm4		NXP PQC [A]		Slower	Smaller
		Runtime	RAM	Runtime	RAM	Runtime	RAM
Dilithium-2	Sign	19 ms	50 kB	61 ms	5 kB	3.2x	10.0x
	Verify	7 ms	11 kB	16 ms	3 kB	2.3x	3.7x
Dilithium-3	Sign	31 ms	69 kB	119 ms	7 kB	3.8x	9.9x
	Verify	12 ms	10 kB	29 ms	3 kB	2.4x	3.3x
Dilithium-5	Sign	42 ms	123 kB	168 ms	8 kB	4.0x	15.4x
	Verify	21 ms	12 kB	50 ms	3 kB	2.4x	4.0x



All Dilithium parameter sets will fit on a device with ~8KB memory.



Price: factor 3 to 4 in performance HW accelerators



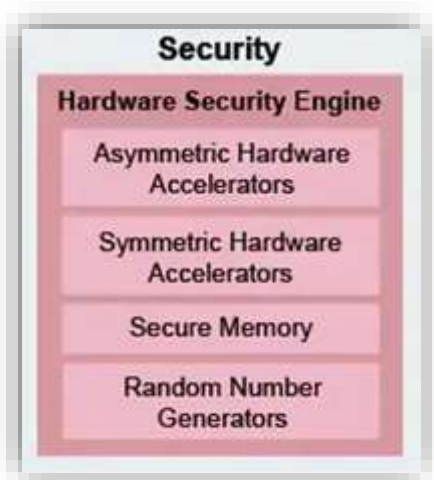


# Example of what we do at NXP

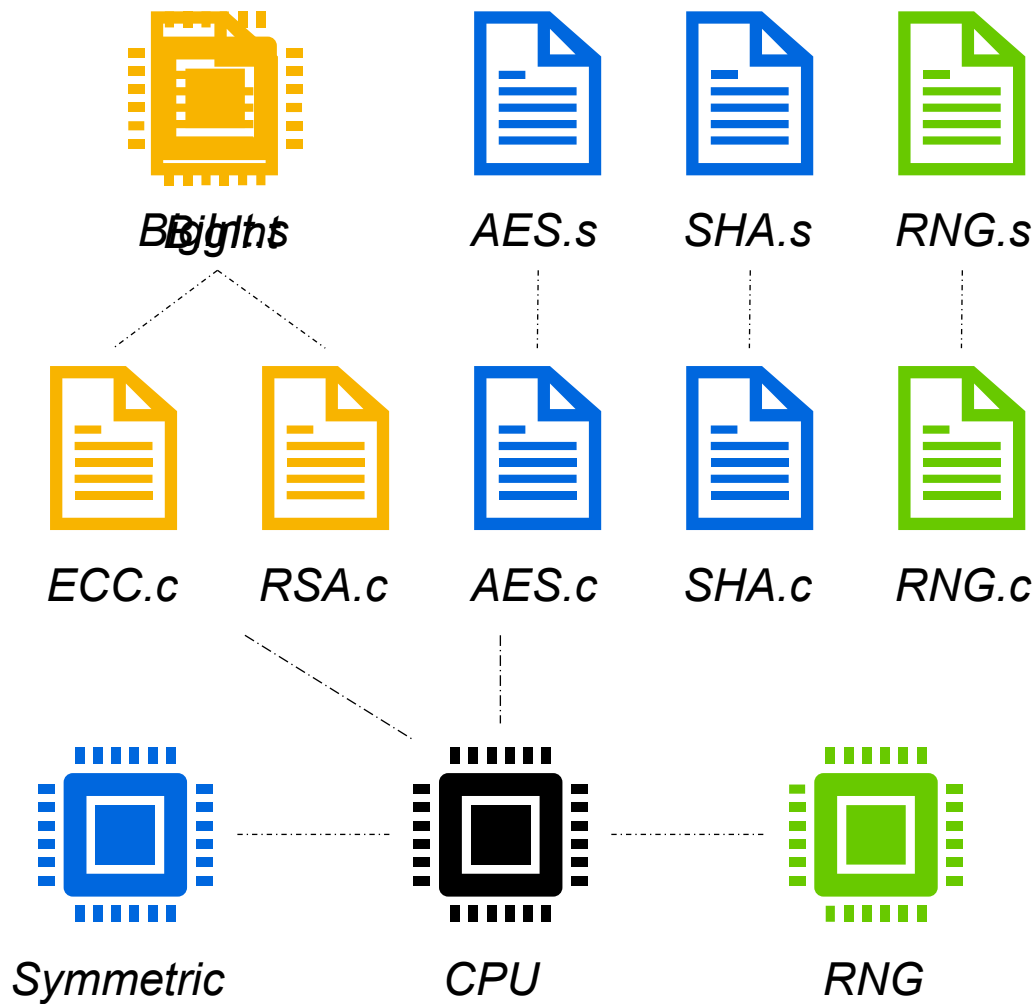
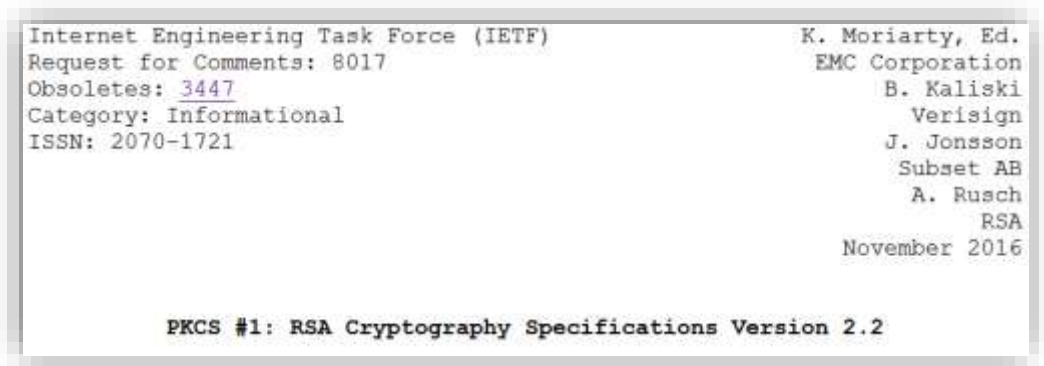
Joppe W. Bos, Joost Renes and Christine van Vredendaal: [\*Polynomial Multiplication with Contemporary Co-Processors: Beyond Kronecker, Schönhage-Strassen & Nussbaumer\*](#). [USENIX Security Symposium 2022](#).



# Implementing Classical cryptography

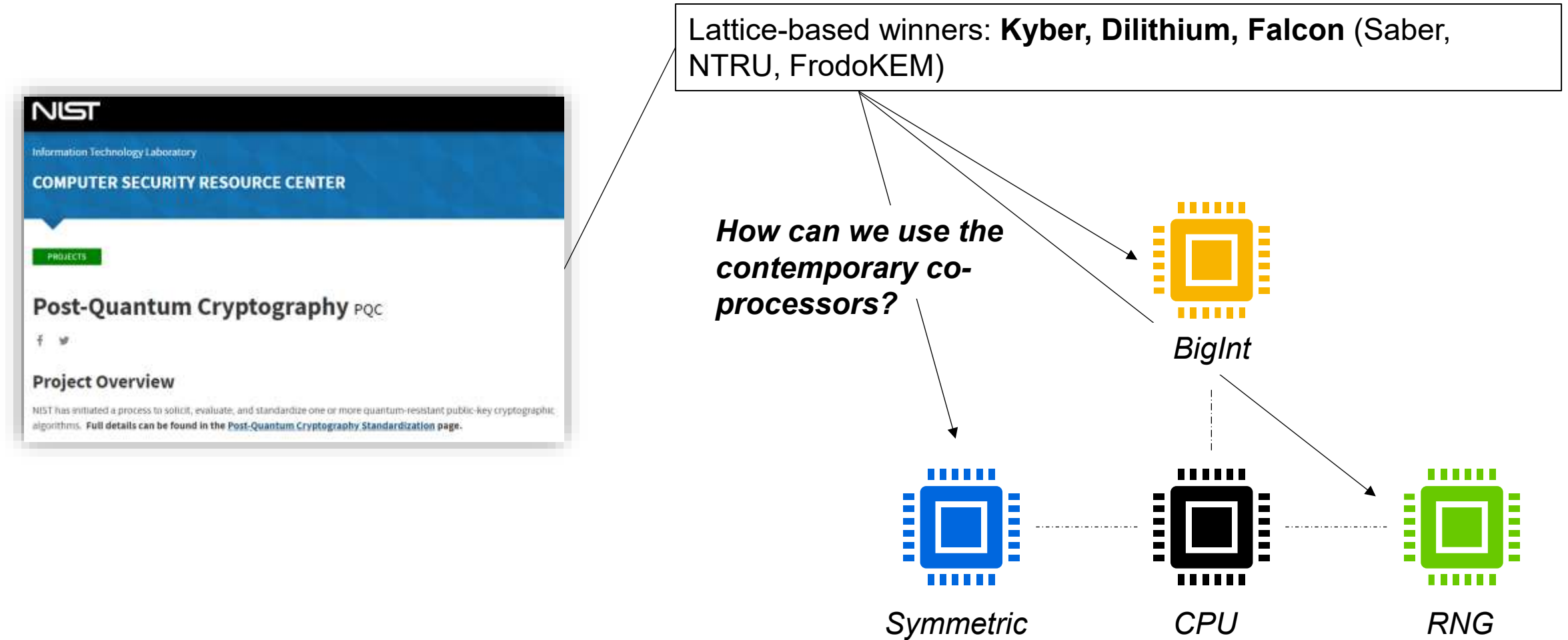


*S32G2 automotive processor spec*





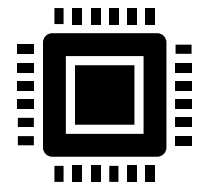
# Implementing post-quantum cryptography





# Re-using existing HW

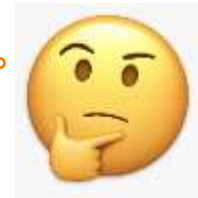
Approach	Core	Structure	Size
RSA	Modular multiplication	$(\mathbb{Z}/n\mathbb{Z})^*$	$n$ is 3072-bit
ECC	Elliptic curve scalar multiplication	$E(\mathbb{F}_p)$	$p$ is 256-bit
Lattice	Polynomial multiplication	$(\mathbb{Z}/q\mathbb{Z})[X]/(X^n + 1)$	$q$ is 16-bit $n$ is 256



Co-pro present in chips



Can we use this?





# Kronecker substitution

*Polynomial domain*

$$f = 1 + 2x + 3x^2 + 4x^3$$

$$g = 5 + 6x + 7x^2 + 8x^3$$

✖

$$fg = \underline{5} + \underline{16x} + \underline{34x^2} + \underline{60x^3} + \underline{61x^4} + \underline{52x^5} + \underline{32x^6}$$

*Kronecker domain (with evaluation point 100)*

$$f(100) = 4030201$$

$$g(100) = 8070605$$

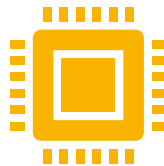
✖

$$fg(100) = \underline{32526160341605}$$

**Grundzüge einer arithmetischen Theorie der  
algebraischen Grössen.**

(Von *L. Kronecker*.)

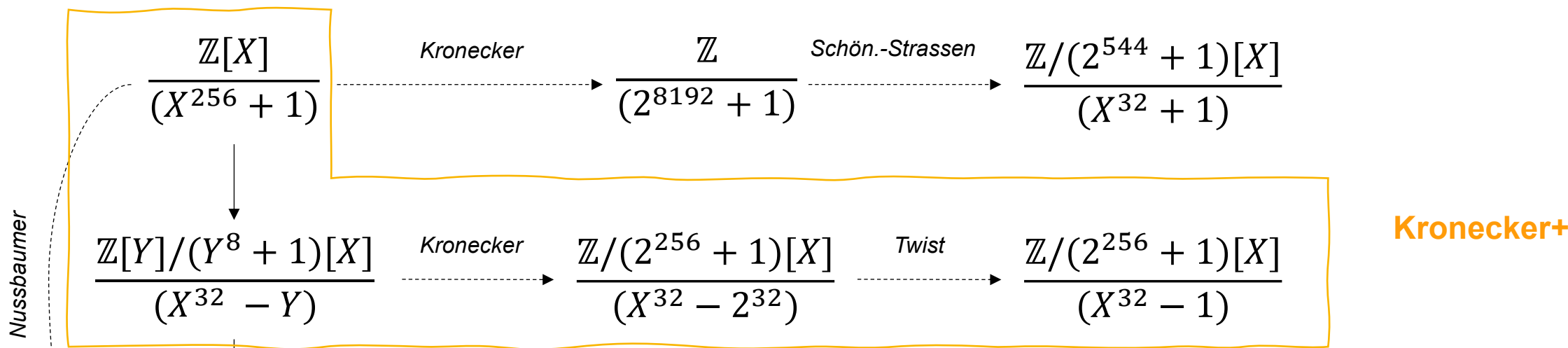
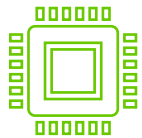
(Abdruck einer Festschrift zu Herrn *E. E. Kummers* Doctor-Jubiläum, 10. September 1881.)





# Polynomial multiplication techniques

Kronecker evaluation at  $2^{32}$   
Multiplication with a **256-bit** multiplier



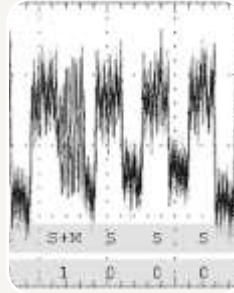
Algorithm	# Muls	# Bits
Kron. + Schoolbook	1024	256
Kron. + Karatsuba	243	256
Kron. + Toom-Cook	63	256
Kron. + Schön.-Strassen	32	544
Nussbaumer + Kron.	64	256
<b>Kronecker+</b>	<b>32</b>	<b>256</b>

[A] Albrecht, Hanser, Hoeller, Pöppelmann, Virdia, Wallner; Implementing RLWE-based schemes using an RSA co-processor. TCHES 2019

[B] Harvey. Faster polynomial multiplication via multipoint Kronecker substitution. J. of Sym. Comp. 2009.



# Resistance against physical & logical attacks



## Side-channel attacks

- Power analysis (SPA, DPA)
- Electromagnetic analysis (SEMA, DEMA)
- Timing Analysis
- Photo-emission microscopy (high-end)
- Profiled, unprofiled and ML-assisted variants



## Fault injection attacks

- Voltage or clock glitching
- Electromagnetic fault injection (EMFI)
- Body bias injection
- Laser fault injection
- Single and multi-shot scenarios

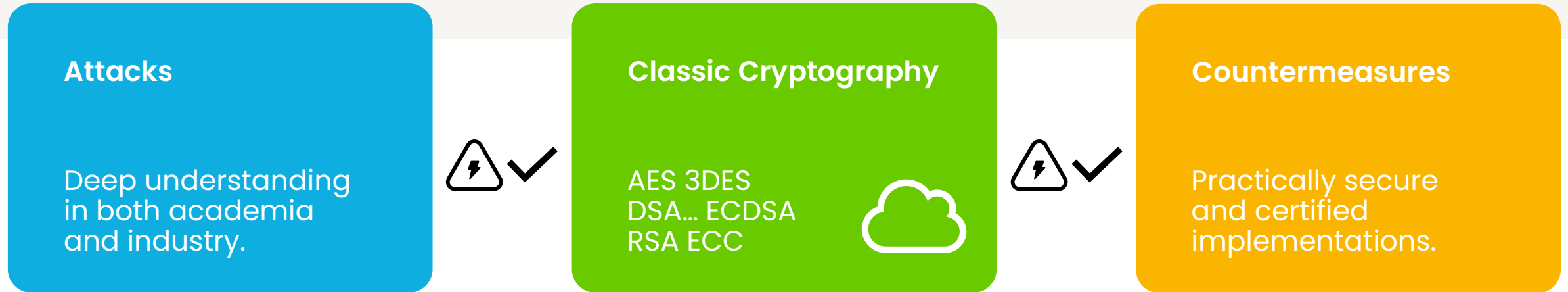


## Invasive attack

- Focused Ion Beam (FIB) modifications
- Micro/Nano-probing of internal signals
- Signal forcing
- Delaying
- Reverse-engineering

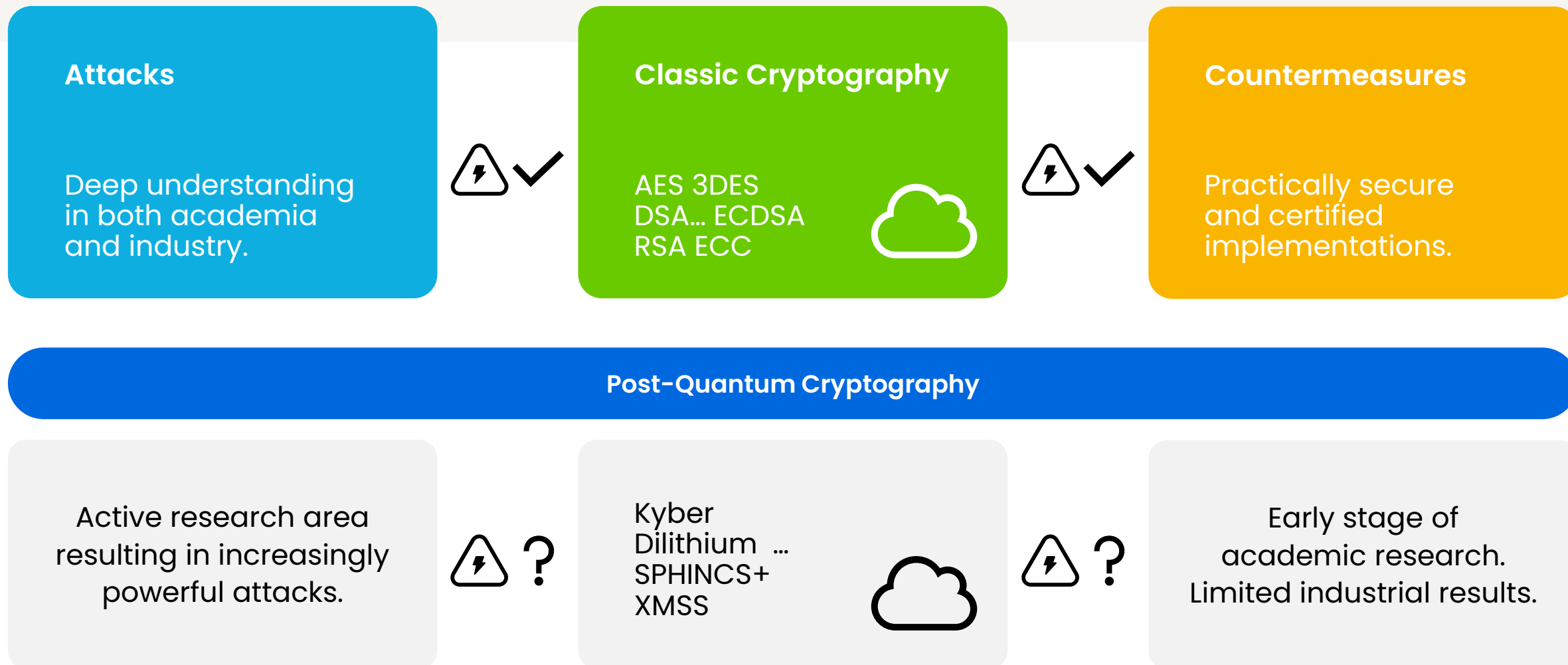


# Embedded cryptography and implementation attacks



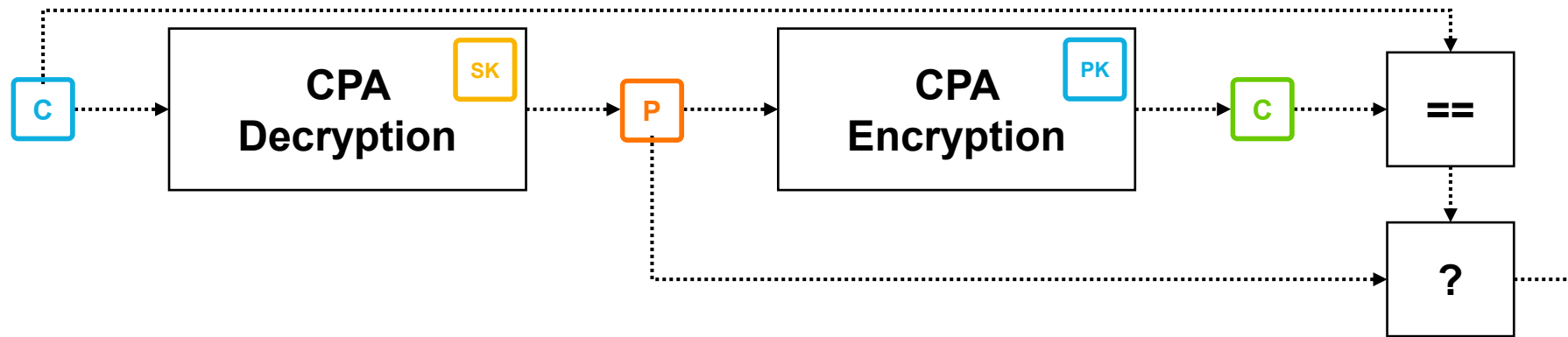


# Embedded cryptography and implementation attacks





# Fujisaki Okamoto transform



Transform a scheme which achieves **IND-CPA**  
("chosen plaintext attack") security to reach **IND-CCA**  
("indistinguishability against chosen-ciphertext attacks") security

Fujisaki, E. and Okamoto  
T., Secure integration of  
asymmetric and symmetric  
encryption schemes, CRYPTO  
1999 and JoC 2013

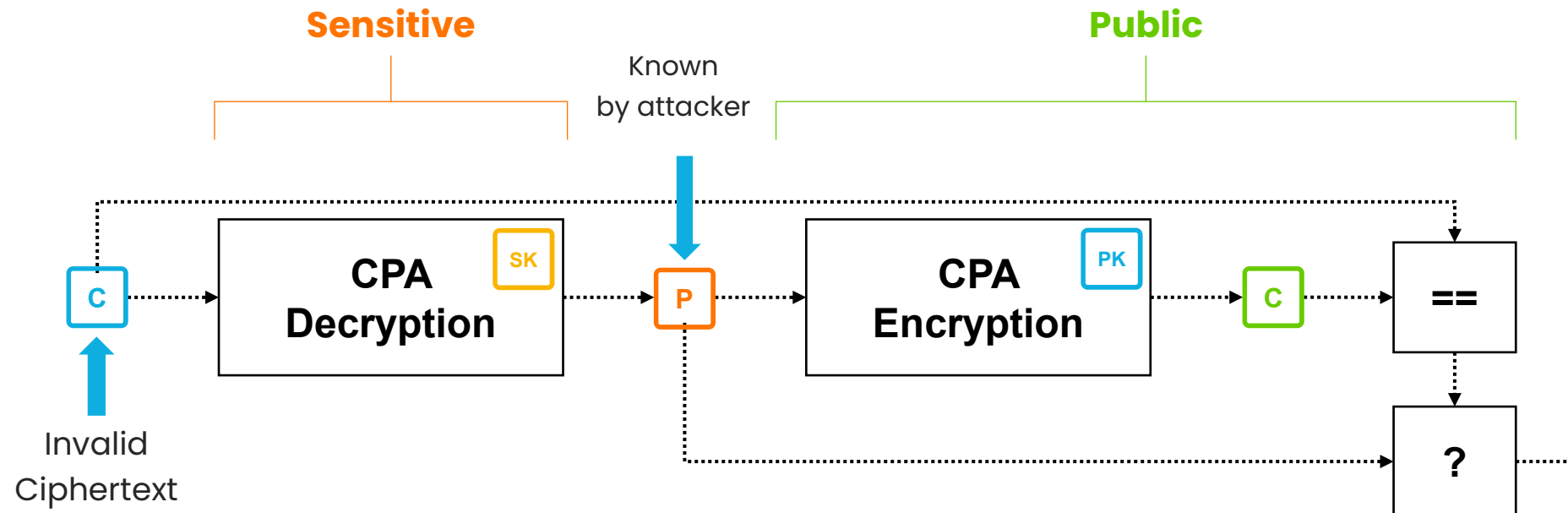


# The SCA Problem of the FO-Transform



## Attack 1: Chosen Plaintext

- Attacker inputs only valid ciphertexts
- Attack focuses on **CPA Decryption**, everything after (and including) **P** is public
- Only need to protect **CPA Decryption**



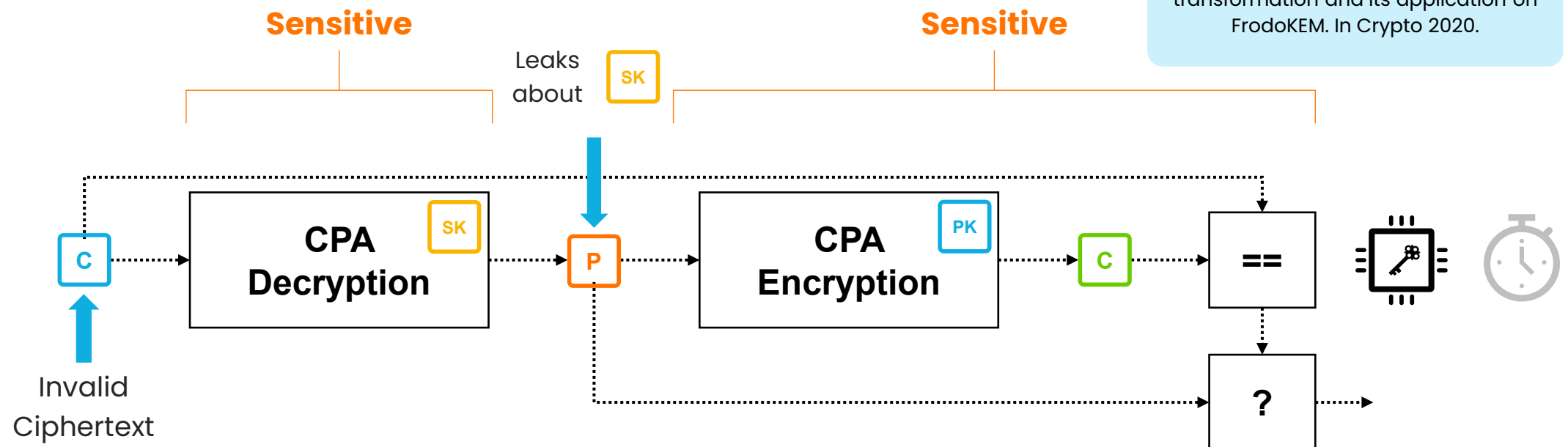


# The SCA Problem of the FO-Transform



## Attack 2: Chosen Ciphertext

- Attacker inputs specially-crafted invalid ciphertexts
- Attack focuses on **CPA Decryption** + everything after (and including) **P** is potentially sensitive
- Potentially all (or most) modules need to be hardened





# From Theory to practice: Secure implementations (NXP PQC Team)

Only with carefully managed maximum number of issued signatures

First completely masked implementation of Kyber / FIPS 203 !

Year	Venue	FIPS 203	FIPS 204	Title
2021	TCHES			Masking Kyber: First- and Higher-Order Implementations
2021	RWC			Post-Quantum Crypto: The Embedded Challenge
2022	TCHES			Post-Quantum Authenticated Encryption against Chosen-Ciphertext SCA
2022	RWC			Surviving the FO-calypse: Securing PQC Implementations in Practice
2023	TCHES			From MLWE to RLWE: A Differential Fault Attack on Randomized & Deterministic Dilithium
2023	TCHES			Protecting Dilithium Against Leakage Revisited Sensitivity Analysis
2024	RWC			Lessons Learning from Protecting CRYSTALS-Dilithium
2024	TCHES			Exploiting Small-Norm Polynomial Multiplication with Physical Attacks
2024	RWC			Challenges of Migration to PQ Secure Embedded Systems

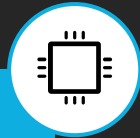
Completely masked implementation of Dilithium / FIPS 204 !



# NXP S32G2 vehicle network processor with PQC integration

## Our target platform: S32G274A

- 3 Lockstep Arm® Cortex®-M7 Microcontrollers
- 4 Cluster Lockstep Cortex-A53 Microprocessors
- 8 MB of System RAM
- Network Accelerators (LLCE/PFE)
- **Hardware Security Engine (HSE)**
- ASIL D Functional Safety Support



## Post-Quantum Crypto

- Integrate PQC secure signature verification
- Protection against Fault Attacks
- Enable PQC secure boot
- Secure Over-the-Air (OTA) updates
- Secure vehicle and driver data

[www.nxp.com/s32g2](http://www.nxp.com/s32g2)





# Benchmarks for authentication of FW signature on the S32G2

Alg.	Size		Performance (ms)			
			1 KB		128 KB	
	PK	Sig.	Inst.	Boot	Inst.	Boot
RSA 4K	512	512	2.6	0.0	2.7	0.2
ECDSA-p256	64	64	6.2	0.0	6.4	0.2
Dilithium-3	1952	3293	16.7	0.0	16.9	0.2



Demonstrator only, further optimizations are possible (such as hardware accelerated SHA-3)



Signature verification only required once for installation!



During boot the signature verification can be replaced with a check of the Reference Proof of Authenticity





# Conclusions

- We are always looking for talented people in math / crypto!
- Need to have an applied interest as well.
- New mathematical techniques to map algorithms to resource constrained devices.
- Software / hardware skills are a plus
- Crypto / number theory knowledge is a must!

Experience shows it is easier to teach software development skills to an applied mathematician than number theory to an engineer 😊

Interested? Job? Internship? Industry PhD with KU Leuven?

Contact me: [joppe.bos@nxp.com](mailto:joppe.bos@nxp.com)







# Get in touch!

**Joppe W. Bos**

joppe.bos@nxp.com

[nxp.com](https://www.nxp.com)





Brighter Together